# Data Representation: bits, bytes, integers 

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## Programming assignments

## Programming assignment 1

- Due date extended to: 11:59pm Monday, February 15.
- Take good advantage of this opportunity.
- Familiarity with C is a vital foundation for this class and future classes.
- Code review discussion during week of February 22 - February 26.


## Programming assignment 2

- Released later today Thursday, February 11.
- Due after two weeks: February 25.
- Same techniques in programming C.
- Review of graph algorithms.


## Looking ahead

## Lecture plan

1. Today, Thursday, 2/11: Data representation of integers.
2. Tuesday, $2 / 16$ : Data representation of floating point numbers.
3. Thursday, 2/18: Data representation of floating point numbers.

Reading assignment

- Computer Systems: A Programmer's Perspective Chapter 2.


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## Why binary

## Everything is bits

- Each bit is 0 or 1
- By encoding/interpreting sets of bits in various ways
- Computers determine what to do (instructions)
- ... and represent and manipulate numbers, sets, strings, etc...

■ Why bits? Electronic Implementation

- Easy to store with bistable elements
- Reliably transmitted on noisy and inaccurate wires



## Why binary

Figure: Rahul Sarpeshkar. Analog Versus Digital: Extrapolating from Electronics to Neurobiology. 1998.


## Why binary

## Digital encodings

Each doubling of either precision or range only needs one additional bit.

## Analog encodings

Each doubling of either precision or range needs doubling of either area or power.

## Decimal, binary, octal, and hexadecimal

| Decimal | Binary | Octal | Hexadecimal |  |  | Decimal | Binary | Octal |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
|  | Hexadecimal |  |  |  |  |  |  |  |
| 0 | 0 b 0000 | 0 o 0 | $0 \times 0$ |  | 8 | 0 b 1000 | 0 o 10 | $0 \times 8$ |
| 1 | 0 b 0001 | 0 o 1 |  | $0 \times 1$ |  | 9 | 0 b 1001 | 0 o 11 |

In C, format specifiers for printf() and fscanf():

1. decimal: '\%d'
2. binary: none
3. octal: ' $\% \mathrm{o}^{\prime}$
4. hexadecimal: ' $\% x^{\prime}$

## Decimal, binary, octal, and hexadecimal

How to represent the range of unsigned char in each?
Unsigned char is one byte, 8 bits.

1. decimal: 0 to 255
2. binary: 0 b 0 to 0 b 11111111
3. octal: 0 to 00377 (group by 3 bits)
4. hexadecimal: $0 \times 00$ to $0 \times \mathrm{FF}$ (group by 4 bits)

## Bitwise operations

Why are bitwise operations important?

- Network and UNIX settings using bit masks (e.g., umask)
- Hardware and microcontroller programming (e.g., Arduinos)
- Instruction set architecture encodings (e.g., ARM, x86)


## Bitwise operations

## ~: bitwise NOT

unsigned char $\mathrm{a}=128$

$$
\begin{aligned}
a & =0 b 1000 \_0000 \\
\sim & =\sim 0 b 1000 \_0000 \\
& =0 b 0111 \_1111 \\
& =127
\end{aligned}
$$

## Bitwise operations

\&: bitwise AND

$$
\begin{aligned}
3 \& 1 & =0 b 11 \& 0 b 01 \\
& =0 b 01 \\
& =1
\end{aligned}
$$

| a | b | $\mathrm{a} \& \mathrm{~b}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 0 |
| 1 | 0 | 0 |
| 1 | 1 | 1 |

## Bitwise operations

I: bitwise OR

$$
\begin{aligned}
3 \mid 1 & =0 b 11 \mid 0 b 01 \\
& =0 b 11 \\
& =3 \\
2 \mid 1 & =0 b 10 \mid 0 b 01 \\
& =0 b 11 \\
& =3
\end{aligned}
$$

| a | b | $\mathrm{a} \mid \mathrm{b}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 1 |

## Bitwise operations

^: bitwise XOR

| a | b | $\mathrm{a}^{\wedge} \mathrm{b}$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

$$
\begin{aligned}
3 \wedge 1 & =0 b 11 \wedge 0 b 01 \\
& =0 b 10 \\
& =2
\end{aligned}
$$

## Don't confuse bitwise operators with logical operators

Bitwise operators

- \&
- 1
- 

Logical operators
-!

- \&\&
- | 1
- != (for bool type)


## Representing characters

USASCII code chart

| $\mathrm{C}_{7} b_{6}$ |  |  |  |  | ${ }^{0}{ }^{0}$ | $0_{0}$ | $\begin{array}{llll}0 & & \\ & 1 & \\ & & 0\end{array}$ | $0^{0} 1$ | ${ }^{1} 00$ | ${ }^{1} 0$ | $110$ | $\begin{array}{llll}1 & & \\ & 1 & \\ & & 1\end{array}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| N/ ${ }^{\text {b }}$ | ${ }^{b_{3}}$ | $\begin{gathered} b_{2} \\ 1 \end{gathered}$ | $\begin{array}{\|c} b_{1} \\ 1 \end{array}$ |  | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| 0 | 0 | 0 | 0 | 0 | NUL | DLE | SP | 0 | 0 | P | , | P |
| 0 | 0 | 0 | 1 | 1 | SOH | DC1 | ! | 1 | A | 0 | 0 | 9 |
| 0 | 0 | 1 | 0 | 2 | STX | DC2 | " | 2 | B | R | b | 1 |
| 0 | 0 | 1 | 1 | 3 | ETX | DC 3 | \# | 3 | C | S | c | \$ |
| 0 | 1 | 0 | 0 | 4 | EOT | DC4 | 1 | 4 | D | $T$ | d | 1 |
| 0 | 1 | 0 | 1 | 5 | ENQ | NAK | \% | 5 | E | U | e | $v$ |
| 0 | 1 | 1 | 0 | 6 | ACK | SYN | 8 | 6 | $F$ | V | 1 | $v$ |
| 0 | 1 | 1 | 1 | 7 | BEL | ETB | , | 7 | G | $w$ | 9 | w |
| 1 | 0 | 0 | 0 | 8 | BS | CAN | 1 | 8 | H | X | n | x |
| 1 | 0 | 0 | 1 | 9 | HT | EM | $)$ | 9 | 1 | Y | $i$ | $y$ |
| 1 | 0 | 1 | 0 | 10 | LF | SUB | * | : | $J$ | Z | 1 | 2 |
| 1 | 0 | 1 | 1 | 11 | VT | ESC | + | ; | K | [ | k | ( |
| 1 | 1 | 0 | 0 | 12 | FF | FS | $\cdots$ | $<$ | L | 1 | 1 | 1 |
| 1 | 1 | 0 | 1 | 13 | CR | GS | - | = | M | ] | m | \} |
| 1 | 1 | 1 | 0 | 14 | SO | RS | . | $>$ | N | $\wedge$ | $n$ | $\sim$ |
| 1 | 1 | 1 | 1 | 15 | SI | US | 1 | ? | 0 | - | 0 | DEL |

Figure: ASCII character set. Image credit Wikimedia

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## Representing negative and signed integers

Ways to represent negative numbers

1. Sign magnitude
2. 1's complement
3. 2's complement

Representing negative and signed integers

Sign magnitude
Flip leading bit.

## Representing negative and signed integers

1's complement

- Flip all bits
- Addition in 1's complement is sound
- In this encoding there are 2 encodings for 0
- $-0: 0 \mathrm{0} 1111$
- +0: 0b0000


## Representing negative and signed integers

2's complement

| signed char | weight in decimal |
| ---: | ---: |
| 00000001 | 1 |
| 00000010 | 2 |
| 00000100 | 4 |
| 00001000 | 8 |
| 00010000 | 16 |
| 00100000 | 32 |
| 01000000 | 64 |
| 10000000 | -128 |

Table: Weight of each bit in a signed char type

- what is the most positive value you can represent? 127
- what is the most negative value you can represent? -128
- how to represent -1? 11111111
- how to represent -2? 11111110


## Representing negative and signed integers

## 2's complement

| signed char | weight in decimal |
| ---: | ---: |
| 00000001 | 1 |
| 00000010 | 2 |
| 00000100 | 4 |
| 00001000 | 8 |
| 00010000 | 16 |
| 00100000 | 32 |
| 01000000 | 64 |
| 10000000 | -128 |

Table: Weight of each bit in a signed char type

- MSB: 1 for negative
- Take the 1's complement number + 1
- Most important; good properties for digital logic


## Importance of paying attention to limits of encoding



Figure: Image credit: CS:APP

## Importance of paying attention to limits of encoding



Figure: Image credit: CS:APP

## Importance of paying attention to limits of encoding



Figure: Image credit: CS:APP

