

Machine-Level Representation of Programs: Bomblab, addressing modes, arithmetic

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Programming Assignment 4: Defusing a Binary Bomb

Unpacking your bomb

Using GDB

2_addressing_modes.s: Understanding source dest operands and memory addressing modes

3_leaq.s: Borrowing memory address calculation to efficiently implement arithmetic

Announcements

Class session plan

- ▶ Today, 3/21: addressing modes (Book chapter 3.4), arithmetic (Book chapter 3.5). Bomblab phase_1.
- ▶ Tuesday, 3/26: Control flow (conditionals, if, for, while, do loops, switch statements) in assembly. (Book chapter 3.6). Bomblab phase_2, phase_3.
- ▶ Thursday, 3/28: Function calls in assembly. (Book chapter 3.7). Bomblab phase_4.
- ▶ Tuesday, 4/2: Arrays and data structures in assembly. (Book chapter 3.8). Bomblab phase_5, phase_6.

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Programming Assignment 4: Defusing a Binary Bomb

Goals

- ▶ Learning to learn to use important tools like GDB.
- ▶ Understand how high level programming constructs compile down to assembly instructions.
- ▶ Practice reverse engineering and debugging.

Setup

- ▶ Programming assignment description PDF on Canvas.
- ▶ Web interface for obtaining bomb and seeing progress.
- ▶ Unpacking.

Unpacking and gathering information about your bomb

What comes in the package

- ▶ `bomb.c`: Skeleton source code
- ▶ `bomb`: The executable binary

```
objdump -t bomb > symbolTable.txt
```

- ▶ `000000000040143a g F .text 0000000000000022 explode_bomb`

```
objdump -d bomb > bomb.s
```

Different phases correspond to different topics about assembly programming in the CS211 lecture slides, in the CS:APP slides, and in the CS:APP book.

- ▶ `phase_1`
- ▶ `phase_2`
- ▶ `explode_bomb`

```
strings -t x bomb > strings.txt
```

Example phase_1 in example bomb from CS:APP website

```
0000000000400ee0 <phase_1>:
```

```
400ee0: 48 83 ec 08          sub     $0x8,%rsp
400ee4: be 00 24 40 00      mov     $0x402400,%esi
400ee9: e8 4a 04 00 00     callq  401338 <strings_not_equal>
400eee: 85 c0              test   %eax,%eax
400ef0: 74 05             je     400ef7 <phase_1+0x17>
400ef2: e8 43 05 00 00     callq  40143a <explode_bomb>
400ef7: 48 83 c4 08       add     $0x8,%rsp
400efb: c3              retq
```

Understanding what we're seeing here

- ▶ Don't let `callq` to `explode_bomb` at instruction address `400ef2` happen...
- ▶ so, must ensure `je` instruction does jump, so we want `test` instruction to set ZF condition code to 0.
- ▶ so, must ensure `callq` to `strings_not_equal()` function returns 0.

Using GDB to carefully step through execution of the bomb program

```
gdb bomb
```

Finding help in GDB

- ▶ `help`: Menu of documentation.
- ▶ `help layout`: Useful tip to use either `layout asm` or `layout regs` for this assignment.
- ▶ `help aliases`
- ▶ `help running`
- ▶ `help data`
- ▶ `help stack`

Using GDB to carefully step through execution of the bomb program

```
gdb bomb
```

Setting breakpoints and running / stepping through code

- ▶ `break explode_bomb` or `b explode_bomb`: Pause execution upon entering `explode_bomb` function.
- ▶ `break phase_1` or `b phase_1`: Pause execution upon entering `phase_1` function.
- ▶ `run mysolution.txt` or `r mysolution.txt`: Run the code passing the solution file.
- ▶ `continue` or `c`: Continue until the next breakpoint.
- ▶ `nexti` or `ni`: Step one instruction, but proceed through subroutine calls.
- ▶ `stepi` or `si`: Step one instruction exactly. Steps into functions / subroutine calls.

Using GDB to carefully step through execution of the bomb program

```
gdb bomb
```

Printing and examining registers and memory addresses

- ▶ `print /x $eax` or `p /x $eax`: Print value of `%eax` register as hex.
- ▶ `print /d $eax` or `p /d $eax`: Print value of `%eax` register as decimal.
- ▶ `x /s 0x402400`: Examine memory address `0x402400` as a string.

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Immediate, register, and memory

Immediate

Constant integer values. Example: `2_addressing_modes.c` `immediate()`

Register

One of the registers of appropriate size for data type. Example: `1_swap.c`

Memory

Access to memory at calculated

movq Operand Combinations

	Source	Dest	Src, Dest	C Analog
movq	Imm	Reg	<code>movq \$0x4, %rax</code>	<code>temp = 0x4;</code>
		Mem	<code>movq \$-147, (%rax)</code>	<code>*p = -147;</code>
	Reg	Reg	<code>movq %rax, %rdx</code>	<code>temp2 = temp1;</code>
		Mem	<code>movq %rax, (%rdx)</code>	<code>*p = temp;</code>
	Mem	Reg	<code>movq (%rax), %rdx</code>	<code>temp = *p;</code>

Cannot do memory-memory transfer with a single instruction

Simple Memory Addressing Modes

■ Normal (R) Mem[Reg[R]]

- Register R specifies memory address
- Aha! Pointer dereferencing in C

```
movq (%rcx), %rax
```

■ Displacement D(R) Mem[Reg[R]+D]

- Register R specifies start of memory region
- Constant displacement D specifies offset

```
movq 8(%rbp), %rdx
```

```
movl 4(%rbp), %edx
```

Normal

Simple pointers.
Example: 2_addressing_modes.c
immediate()

Displacement

Array access with constant index.
Example: 2_addressing_modes.c
displacement()

Complete Memory Addressing Modes

■ Most General Form

$D(Rb, Ri, S)$ $Mem[Reg[Rb]+S*Reg[Ri]+ D]$

- D: Constant “displacement” 1, 2, or 4 bytes
- Rb: Base register: Any of 16 integer registers
- Ri: Index register: Any, except for `%rsp`
- S: Scale: 1, 2, 4, or 8 (*why these numbers?*)

■ Special Cases

(Rb, Ri) $Mem[Reg[Rb]+Reg[Ri]]$

$D(Rb, Ri)$ $Mem[Reg[Rb]+Reg[Ri]+D]$

(Rb, Ri, S) $Mem[Reg[Rb]+S*Reg[Ri]]$

Indexed

Array access with
variable index.

Example: `2_addressing_modes.c`
`index()`

Address Computation Examples

<code>%rdx</code>	<code>0xf000</code>
<code>%rcx</code>	<code>0x0100</code>

Expression	Address Computation	Address
<code>0x8(%rdx)</code>	<code>0xf000 + 0x8</code>	<code>0xf008</code>
<code>(%rdx,%rcx)</code>	<code>0xf000 + 0x100</code>	<code>0xf100</code>
<code>(%rdx,%rcx,4)</code>	<code>0xf000 + 4*0x100</code>	<code>0xf400</code>
<code>0x80(,%rdx,2)</code>	<code>2*0xf000 + 0x80</code>	<code>0x1e080</code>

2_addressing_modes.c: Imm→Mem

C code

```
void immediate ( long * ptr ) {  
    *ptr = 0xFFFFFFFFFFFFFFFF;  
}
```

Assembly code

```
immediate:  
    movq $-1, (%rdi)  
    ret
```

- ▶ \$ indicates the immediate value; corresponds to literals in C
- ▶ (%rdi) indicates memory location at address stored in %rdi register

2_addressing_modes.c: Imm→Mem (with displacement)

C code

```
void displacement_l ( long * ptr ) {  
    ptr[1] = 0xFFFFFFFFFFFFFFFF;  
}
```

- ▶ `8(%rdi)` indicates memory location at address stored in `%rdi` register + 8

Assembly code

```
displacement_l:  
    movq $-1, 8(%rdi)  
    ret
```

2_addressing_modes.c: Imm→Mem (with displacement)

function signature	assembly code
<code>void displacement_c (char * ptr);</code>	<code>movb \$-1, 1(%rdi)</code>
<code>void displacement_s (short * ptr);</code>	<code>movw \$-1, 2(%rdi)</code>
<code>void displacement_i (int * ptr);</code>	<code>movl \$-1, 4(%rdi)</code>
<code>void displacement_l (long * ptr);</code>	<code>movq \$-1, 8(%rdi)</code>

2_addressing_modes.c: Imm→Mem (with index)

C code

```
void index_l ( long * ptr, long index ) {  
    ptr[index] = 0xFFFFFFFFFFFFFFFF;  
}
```

- ▶ `(%rdi,%rsi,8)` indicates memory location at address stored in `%rdi` register + $8 \times$ value stored in `%rsi` register

Assembly code

```
index_l:  
    movq $-1, (%rdi,%rsi,8)  
    ret
```

2_addressing_modes.c: Imm→Mem (with index)

function signature	assembly code
<code>void index_c (char * ptr, long index);</code>	<code>movb \$-1, (%rdi,%rsi)</code>
<code>void index_s (short * ptr, long index);</code>	<code>movw \$-1, (%rdi,%rsi,2)</code>
<code>void index_i (int * ptr, long index);</code>	<code>movl \$-1, (%rdi,%rsi,4)</code>
<code>void index_l (long * ptr, long index);</code>	<code>movq \$-1, (%rdi,%rsi,8)</code>

2_addressing_modes.c: Imm→Mem (with displacement and index)

C code

```
void displacement_and_index ( long * ptr, long index ) {  
    ptr[index+1] = 0xFFFFFFFFFFFFFFFF;  
}
```

- ▶ `8(%rdi,%rsi,8)` indicates memory location at address stored in `%rdi` register + $8 \times$ value stored in `%rsi` register + 8

Assembly code

```
displacement_and_index:  
    movq $-1, 8(%rdi,%rsi,8)  
    ret
```

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Address Computation Instruction

■ `leaq Src, Dst`

- Src is address mode expression
- Set Dst to address denoted by expression

■ Uses

- Computing addresses without a memory reference
 - E.g., translation of `p = &x[i];`
- Computing arithmetic expressions of the form $x + k*y$
 - $k = 1, 2, 4, \text{ or } 8$

■ Example

```
long m12(long x)
{
    return x*12;
}
```

Converted to ASM by compiler:

```
leaq (%rdi,%rdi,2), %rax # t <- x+x*2
salq $2, %rax           # return t<<2
```

Example: 3_leaq.c

Load effective address

```
1 long * leaq (  
2     long * ptr, long index  
3 ) {  
4     return &ptr[index+1];  
5 }
```

```
1 long mulAdd (  
2     long base, long index  
3 ) {  
4     return base+index*8+8;  
5 }
```

Both C code functions above translate to the assembly on the right.

leaq:

mulAdd:

```
leaq 8(%rdi,%rsi,8), %rax  
ret
```

Explanation

- ▶ `leaq src, dest` takes the effective address of the memory (index, displacement) expression of `src` and puts it in `dest`.
- ▶ `leaq` has shorter latency (takes fewer CPU cycles) than `imulq`, so GCC will use `leaq` whenever it can to calculate expressions like $y + ax + b$.