

# Languages and representations for quantum computing: Stabilizer formalism

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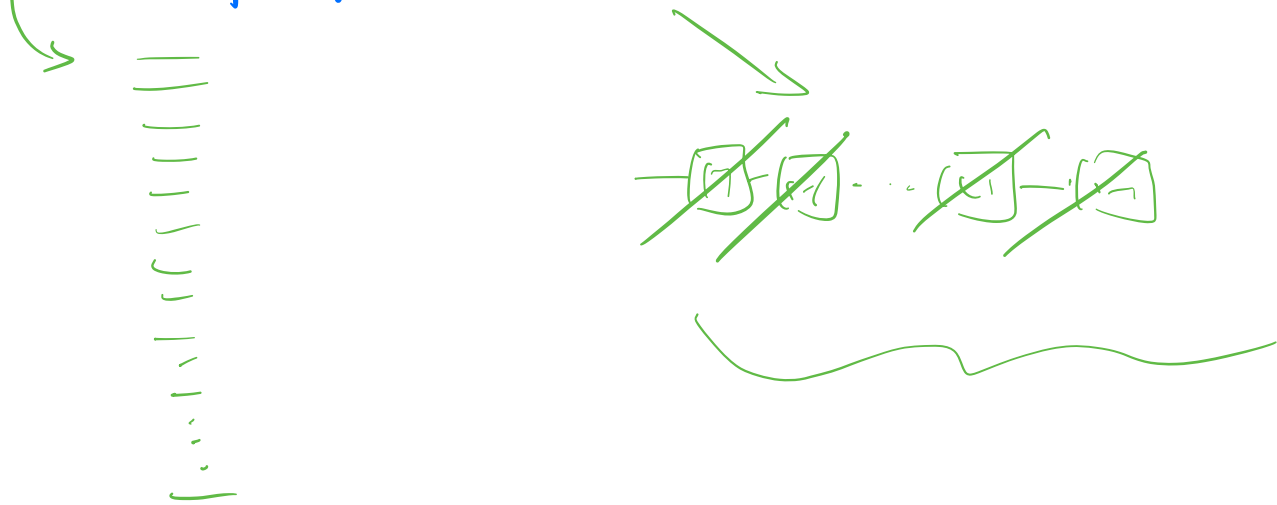
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# Quantum Circuit Simulation

What makes a circuit hard to simulate?

# of qubits? (width)

# of operations? (depth)



large superpositions?



high fidelity?

large entanglement?

# non Cliffords

tree width,

What is it that gives quantum computers an advantage compared to classical computing?

- ▶ Superposition?
- ▶ Entanglement?
- ▶ Both?
- ▶ Neither?

# Importance of representations in quantum intuition, programming, and simulation

- ▶ Conventional quantum circuits and state vector view of QC conceals symmetries, hinders intuition.
- ▶ Classical simulation of quantum computing is actually tractable for *a certain subset* of quantum gates.
- ▶ Both the logical and native gatesets in a quantum architecture need to be *universal* for quantum advantage.

# Several views/representations of quantum computing

- ▶ Programming has several views: functional programming, procedural programming.
- ▶ Physics has several views: Newtonian, Lagrangian, Hamiltonian

Different views reveal different symmetries, offer different intuition.

# Several views/representations of quantum computing

- ▶ Schrödinger: state vectors and density matrices
- ▶ Heisenberg: stabilizer formalism
- ▶ Tensor-network
- ▶ Feynman: path sums

A survey of these representations of quantum computing is given in Chapter 9 of this recent book [Ding and Chong, 2020].

# Several views/representations of quantum computing

- ▶ Schrödinger: state vectors and density matrices
- ▶ Heisenberg: stabilizer formalism
- ▶ Tensor-network
- ▶ Feynman: path sums
- ▶ Binary decision diagrams (new?)
- ▶ Logical satisfiability equations





# Heisenberg view / stabilizer formalism

- ▶ In Heisenberg quantum mechanics description, emphasis on how *operators* evolve.
- ▶ If we limit operations to the Clifford gates (a subset of quantum gates), simulation tractable in polynomial time and space.
- ▶ Covers some quantum algorithms: quantum superdense coding, quantum teleportation, Deutsch-Jozsa, Bernstein-Vazirani, quantum error correction, most quantum error correction protocols.
- ▶ A model for probabilistic (but not quantum) computation.

# Concrete example on Bell state circuit

$$CNOT_{0,1}(H_0 \otimes I_1) |00\rangle$$

1. Start with N qubits with initial state  $|0\rangle^{\otimes N}$ .
2. Represent the state as its group of stabilizers— $|00\rangle : \{IZ, ZI\}$
3. When simulating the quantum circuit, decompose the Clifford gates to stabilizer gates  $\{CNOT, H, P\}$ .
4. Apply each of the stabilizer gates to the stabilizer representation.
  - ▶ Hadamard on first qubit— $|+\rangle |0\rangle : \{IZ, XI\}$
  - ▶ CNOT on both qubits— $\frac{|00\rangle + |11\rangle}{\sqrt{2}} : \{ZZ, XX\}$

# Representing a state as its group of stabilizers

- ▶ A unitary operator  $U$  stabilizes a pure state  $|\psi\rangle$  if  $U|\psi\rangle = |\psi\rangle$

# Representing a state as its group of stabilizers

1.  $I$  stabilizes everything.

2.  $-I$  stabilizes nothing.

3.  $X$  stabilizes  $|+\rangle$ :  $X|+\rangle = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{bmatrix} = |+\rangle$

4.  $-X$  stabilizes  $|-\rangle$ :  $-X|-\rangle = -\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{-1}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{-1}{\sqrt{2}} \end{bmatrix} = |-\rangle$

5.  $Y$  stabilizes  $|+i\rangle$ :  $Y|+i\rangle = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} \end{bmatrix} = |+i\rangle$

6.  $-Y$  stabilizes  $|-i\rangle$ :  $-Y|-i\rangle = -\begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{-i}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{-i}{\sqrt{2}} \end{bmatrix} = |-i\rangle$

7.  $Z$  stabilizes  $|0\rangle$ :  $Z|0\rangle = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \end{bmatrix} = |0\rangle$

8.  $-Z$  stabilizes  $|1\rangle$ :  $-Z|1\rangle = -\begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} 0 \\ 1 \end{bmatrix} = |1\rangle$

# Representing a state as its group of stabilizers

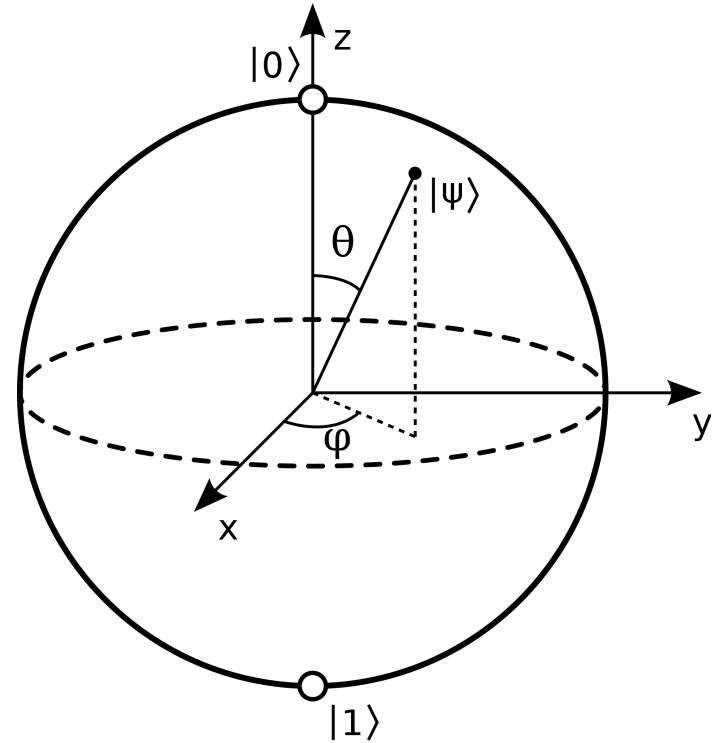
In other words,

1.  $|0\rangle$  is stabilized by  $\{I, Z\}$
2.  $|1\rangle$  is stabilized by  $\{I, -Z\}$
3.  $|+\rangle$  is stabilized by  $\{I, X\}$
4.  $|-\rangle$  is stabilized by  $\{I, -X\}$
5.  $|+i\rangle$  is stabilized by  $\{I, Y\}$
6.  $|-i\rangle$  is stabilized by  $\{I, -Y\}$

# Special places on the Bloch sphere

$$\begin{aligned} |\psi\rangle &= \alpha |0\rangle + \beta |1\rangle \\ &= |\alpha| [\cos(\gamma) + i \cdot \sin(\gamma)] |0\rangle \\ &\quad + |\beta| [\cos(\gamma + \phi) + i \cdot \sin(\gamma + \phi)] |1\rangle \\ &= \cos\left(\frac{\theta}{2}\right) e^{i\gamma} |0\rangle + \sin\left(\frac{\theta}{2}\right) e^{i(\gamma+\phi)} |1\rangle \end{aligned}$$

Enforces  $|\alpha|^2 + |\beta|^2 = 1$

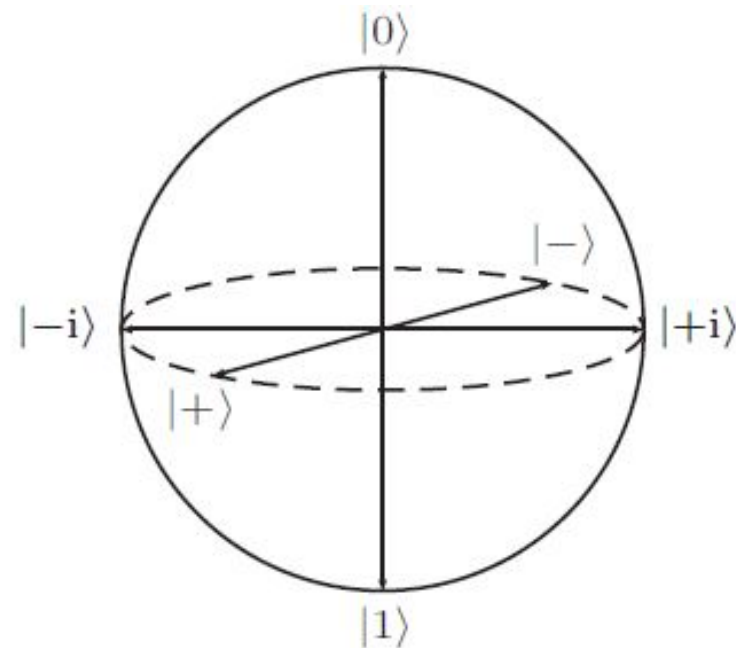


**Figure:** Bloch sphere showing pole states.  
Source: Wikimedia.

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Enforces  $|\alpha|^2 + |\beta|^2 = 1$



**Figure:** Bloch sphere showing pole states.  
Source: Wikimedia.



# Representing a state as its group of stabilizers

For multi-qubit states, the group of stabilizers is the cartesian product of the single-qubit stabilizers

- ▶  $|00\rangle = |0\rangle \otimes |0\rangle$  is stabilized by  $\{I \otimes I, I \otimes Z, Z \otimes I, Z \otimes Z\}$
- ▶  $\frac{|00\rangle + |10\rangle}{\sqrt{2}} = \frac{|0\rangle + |1\rangle}{\sqrt{2}} \otimes |0\rangle = |+\rangle \otimes |0\rangle$  is stabilized by  $\{I \otimes I, I \otimes Z, X \otimes I, X \otimes Z\}$
- ▶  $\frac{|00\rangle + |11\rangle}{\sqrt{2}}$  is stabilized by  $\{I \otimes I, X \otimes X, -Y \otimes Y, Z \otimes Z\}$

# Representing a state as its group of stabilizers

- ▶ Critical result from group theory: for any  $N$ -qubit stabilized state, only  $N$  elements needed to specify group—a result from abstract algebra group theory [Nielsen and Chuang, 2002, Appendix 2]
- ▶ So long as the quantum circuit consists only of Clifford gates, only  $N$  elements needed to specify whole quantum state.
- ▶ Contrast against  $2^N$  amplitudes needed to specify a general  $N$ -qubit quantum state vector.
- ▶ For example a two-qubit states needs four amplitues  $\{a_0, a_1, a_2, a_3\}$  to specify quantum state  $|\psi\rangle = a_0 |00\rangle + a_1 |01\rangle + a_2 |10\rangle + a_3 |11\rangle$ .

# Representing a state as its group of stabilizers

Critical result from group theory: for any N-qubit stabilized state, only N elements needed to specify group.

1.  $|0\rangle$  is stabilized by  $\{I, Z\}$ ,  $Z$  is generator
2.  $|1\rangle$  is stabilized by  $\{I, -Z\}$ ,  $-Z$  is generator
3.  $|+\rangle$  is stabilized by  $\{I, X\}$ ,  $X$  is generator
4.  $|-\rangle$  is stabilized by  $\{I, -X\}$ ,  $-X$  is generator
5.  $|+i\rangle$  is stabilized by  $\{I, Y\}$ ,  $Y$  is generator
6.  $|-i\rangle$  is stabilized by  $\{I, -Y\}$ ,  $-Y$  is generator
7.  $|0\rangle \otimes |0\rangle$  is stabilized by  $\{I \otimes I, I \otimes Z, Z \otimes I, Z \otimes Z\}$ ,  $\{I \otimes Z, Z \otimes I\}$  is generator
8.  $|+\rangle \otimes |0\rangle$  is stabilized by  $\{I \otimes I, I \otimes Z, X \otimes I, X \otimes Z\}$ ,  $\{I \otimes Z, X \otimes I\}$  is generator
9.  $\frac{|00\rangle + |11\rangle}{\sqrt{2}}$  is stabilized by  $\{I \otimes I, X \otimes X, -Y \otimes Y, Z \otimes Z\}$ ,  $\{X \otimes X, Z \otimes Z\}$  is generator

# Concrete example on Bell state circuit

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# Stabilizer gates: $\{CNOT, H, P\}$

1. Hadamard gate: induces superpositions.

2. CNOT gate: induces entanglement.

3. Phase gate: induces complex phases.  $P = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix}$

▶ Despite featuring superposition, entanglement, and complex amplitudes, is *not* universal for quantum computing.

▶ We shall see that the deeply symmetrical structure of these gates prevent access to full quantum Hilbert space.

# Stabilizer gates are a generator for Pauli gates (i.e., Clifford gates decompose to stabilizer gates)

Pauli gates are rotations around respective axes by  $\pi$ .

$$\blacktriangleright Z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} = PP$$

$$\blacktriangleright X = \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{-1}{\sqrt{2}} \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{-1}{\sqrt{2}} \end{bmatrix} = HZH$$

$$\blacktriangleright Y = iXZ$$

$$\blacktriangleright X^2 = Y^2 = Z^2 = I$$

▶ Symmetry is similar to quaternions.

▶ With Clifford gates consisting of  $\{CNOT, H, P, I, X, Y, Z\}$ , sufficient to build many quantum algorithms, including: quantum superdense coding, quantum teleportation, Deutsch-Jozsa, Bernstein-Vazirani, quantum error correction, most quantum error correction protocols.

# Single qubit stabilizer gates bounce stabilizer states around an octahedron on the Bloch sphere

$$P|0\rangle = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} 1 \\ 0 \end{bmatrix} = |0\rangle$$

$$P|1\rangle = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} 0 \\ 1 \end{bmatrix} = i|1\rangle$$

$$P|+\rangle = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} \end{bmatrix} = |+i\rangle$$

$$P|-\rangle = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{-1}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{-i}{\sqrt{2}} \end{bmatrix} = |-i\rangle$$

$$P|+i\rangle = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{-1}{\sqrt{2}} \end{bmatrix} = |-\rangle$$

$$P|-i\rangle = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{-i}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{bmatrix} = |+\rangle$$

# Single qubit stabilizer gates bounce stabilizer states around an octahedron on the Bloch sphere

$$H|0\rangle = \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{-1}{\sqrt{2}} \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{bmatrix} = |+\rangle$$

$$H|1\rangle = \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{-1}{\sqrt{2}} \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{-1}{\sqrt{2}} \end{bmatrix} = |-\rangle$$

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$$H|-i\rangle = \begin{bmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & \frac{-1}{\sqrt{2}} \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{-i}{\sqrt{2}} \end{bmatrix} = \begin{bmatrix} \frac{1}{\sqrt{2}} \\ \frac{i}{\sqrt{2}} \end{bmatrix} = |+i\rangle$$



# Apply each of the stabilizer gates to the stabilizer representation.

## ► Phase:

1.  $Z \rightarrow Z$
2.  $-Z \rightarrow -Z$
3.  $X \rightarrow Y$
4.  $-X \rightarrow -Y$
5.  $Y \rightarrow -X$
6.  $-Y \rightarrow X$

## ► Hadamard:

1.  $Z \rightarrow X$
2.  $-Z \rightarrow -X$
3.  $X \rightarrow Z$
4.  $-X \rightarrow -Z$
5.  $Y \rightarrow -Y$
6.  $-Y \rightarrow Y$

# Single qubit stabilizer gates bounce stabilizer states around an octahedron on the Bloch sphere

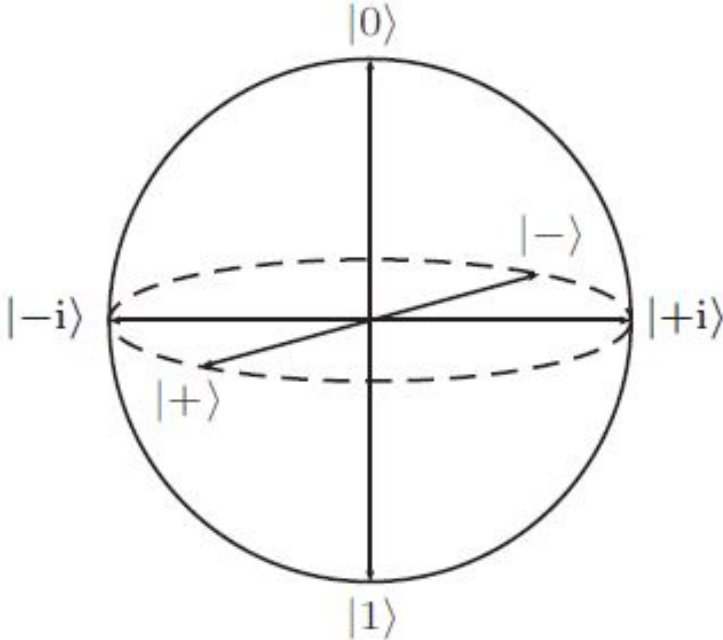


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# Gottesman-Knill theorem and its implications

- ▶ Gottesman-Knill theorem states that there exists a classical algorithm that simulates any stabilizer circuit in polynomial time.
- ▶ Any quantum state created by a Clifford circuit, even if it has lots of superpositions and entanglement, is easy to classically simulate.
- ▶ Quantum computers need at least one non-Clifford gate to achieve universal quantum computation.
- ▶ The T gate, where  $TT = P$ ,  $PP = Z$  is one common choice.
- ▶ There are results showing that a quantum circuit is only exponentially hard to simulate w.r.t. the number of T-gates.

# References

- ▶ Main sources: [Gottesman, 1998] [Aaronson, ]
- ▶ Further reference on separation of probabilistic and quantum computing: [Van Den Nes, 2010]
- ▶ Further reference on applications in classical simulation of Clifford quantum circuits: [Aaronson and Gottesman, 2004]
- ▶ Further reference on applications in classical simulation of general quantum circuits: [Bravyi and Gosset, 2016]

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